EECS3311 Software Design (Fall 2020)

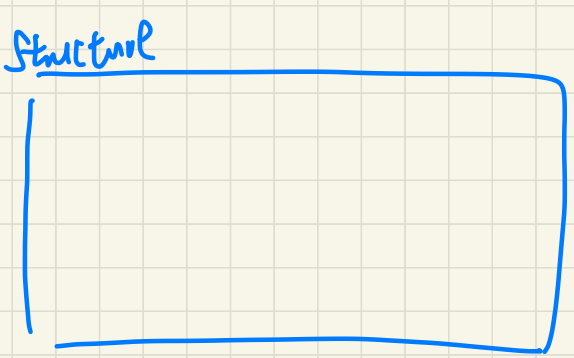
Q&A - Lecture Series W11

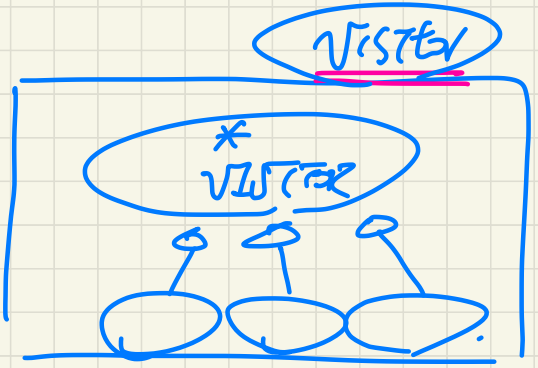
Monday, November 30

The name of visitor comes from the fact that we want to traverse through all the items in a data structure?

recursive.

Can we review this pattern again because I cannot distinguish it from composite pattern?





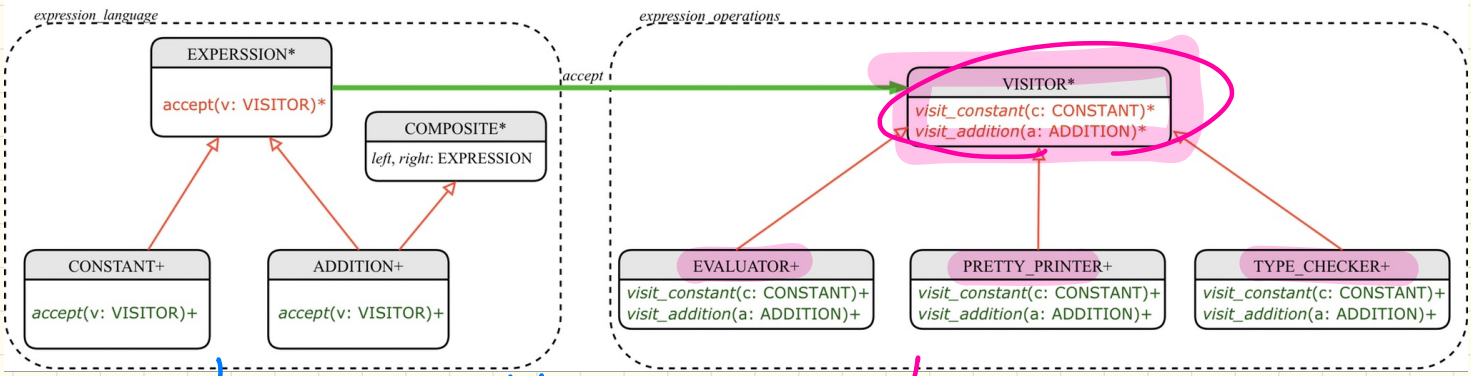
Use visitor if

(recursive)

① already composite implemented

② multiple operations to support

Visitor Design Pattern: Architecture



↳ Composite Cability to build recursive objects)

How to Use Visitors

```
1  test_expression_evaluation: BOOLEAN
2   local add, c1, c2: EXPRESSION ; v: VISITOR
3   do
4    create {CONSTANT} c1.make (1) ; create {CONSTANT} c2.make (2)
5    create {ADDITION} add.make (c1, c2)
6    create {EVALUATOR} v.make
7    add.accept (v)
8    check attached {EVALUATOR} v as eval then
9     Result := eval.value = 3
10   end
11  end
```

v. value ✗

→ using Composite

using Visitor.

I would like to verify that if I want to implement a MULTIPLICATION class,
I must be aware that I will need it ,
and I should have implemented it in the beginning along with the ADD class
instead of adding the MULTIPLICATION class in after everything is implemented.
Is that correct?

① the "structures" cluster in the first working version of the system should contain as many classes as you can anticipate.

② Otherwise, each addition of a new structure class (e.g. mult.) will violate SCP.

context-free grammar

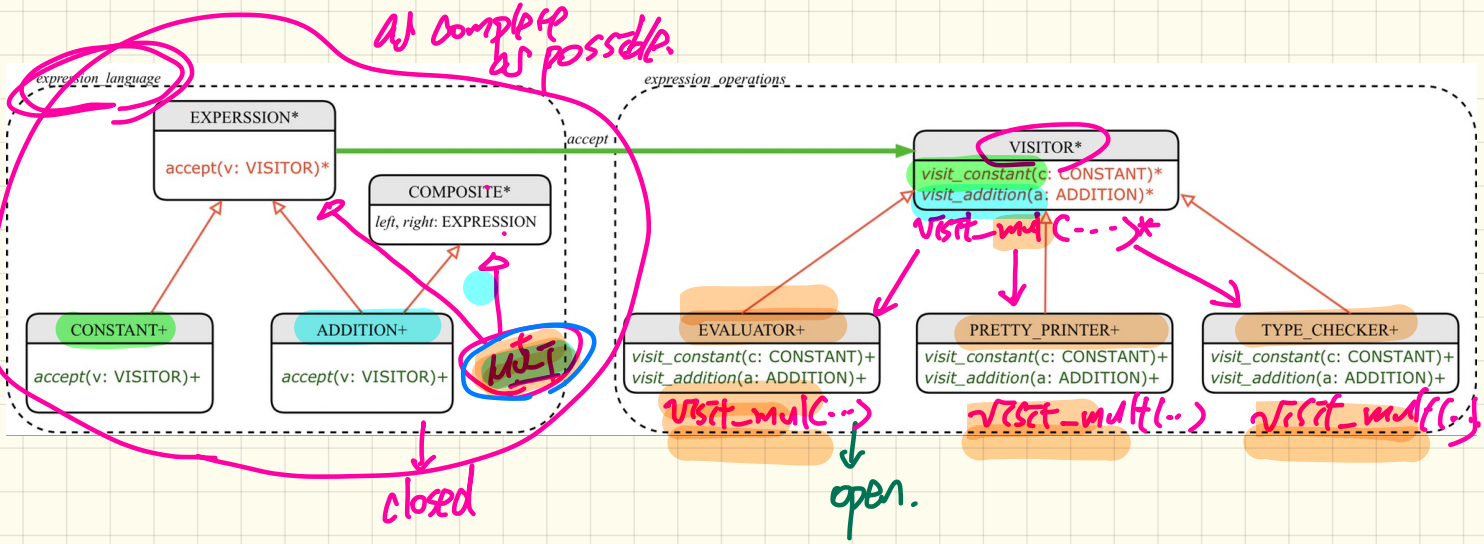
Expression	::=	*IntegerConstant*
	|	*BooleanConstant*
	|	(*BinaryOp*)
	|	(*UnaryOp*)
	|	*CallChain*
IntegerConstant	::=	(1 | 2 | 3 | 4 | 5 | 6 | 7 | 8 | 9)(0 | 1 | 2 | 3 | 4 | 5 | 6 | 7 | 8 | 9)*
BooleanConstant	::=	True
	|	False
BinaryOp	::=	*Expression* + *Expression*
	|	*Expression* - *Expression*
	|	*Expression* * *Expression*
	|	*Expression* / *Expression*
	|	*Expression* % *Expression*
	|	*Expression* && *Expression*
	|	*Expression* || *Expression*
	|	*Expression* == *Expression*
	|	*Expression* > *Expression*
	|	*Expression* < *Expression*
UnaryOp	::=	- *Expression*
	|	! *Expression*
CallChain	::=	*Name*(.*Name*)*

there is a correlation between classes in expression cluster (constant, addition) to the commands of expression operations cluster (visit_constant, visit_addition) while the expression language is closed and we are not going to make changes, then being open in the expression operation cluster means what?

↳ ① add as many descendant classes of VISITOR as you like

② Each new descendant class you add, that's the single point of changes.

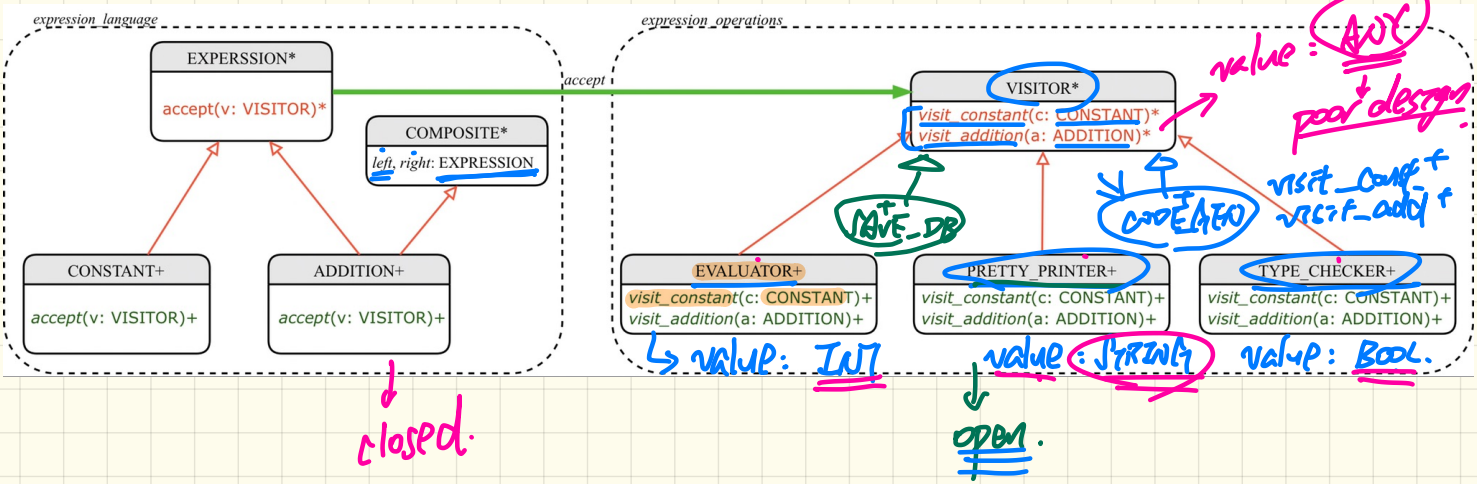
Visitor Pattern: Open-Closed and Single-Choice Principles



What if a new language construct is added?

If the visitor pattern is adopted, what should be closed?

Visitor Pattern: Open-Closed and Single-Choice Principles



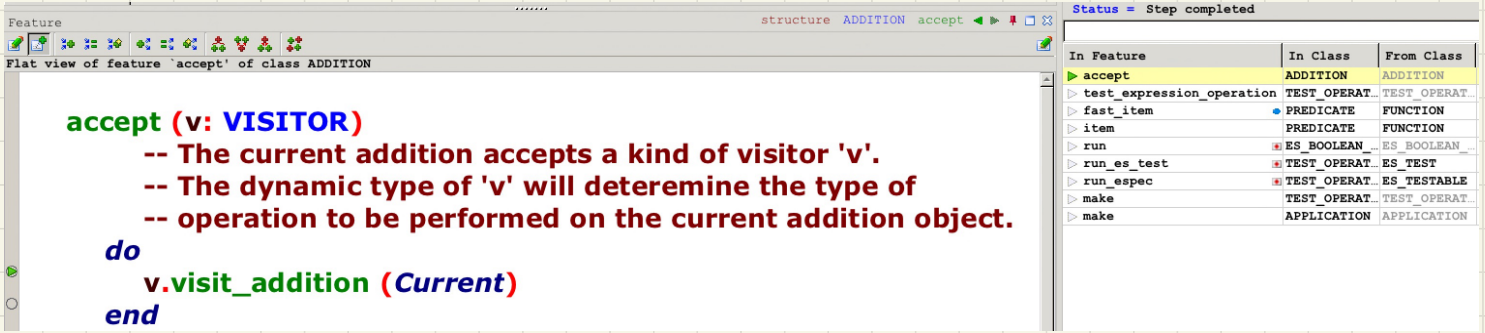
What if a new language operation is added?

↳ SCP satisfied

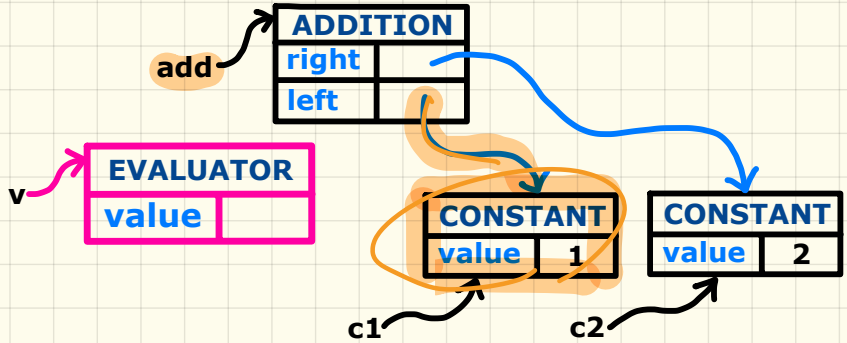
If the visitor pattern is adopted, what should be open?

From Composite-Visitor Totorial:

I can tell the static type of v is VISITOR by "(v: VISITOR)",
but why is its dynamic type is EVALUATOR?
I couldn't find a line code to assign that.
Is it because EVALUATOR is the only class that has the "visit_constant" effective, so EVALUATOR automatically becomes the default type when "visit_constant" is called.
What if we have the PRETTY_PRINTER class, how can the second dispatch decide which one will it call the "visit_constant".



Executing Composite and Visitor Patterns at Runtime



Tracing add.accept(v)
Double Dispatch

1st D.S.

exp.accept(visitor)
↳ triggers a D.S.

```
deferred class VISITOR
  visit_constant(c: CONSTANT) deferred end
  visit_addition(a: ADDITION) deferred end
end
```

```
class EVALUATOR inherit VISITOR
  value: INTEGER
  visit_constant(c: CONSTANT)  do value := c.value end
  visit_addition(a: ADDITION)
    local eval_left, eval_right: EVALUATOR
    do a.left.accept(eval_left)
       a.right.accept(eval_right)
       value := eval_left.value + eval_right.value
    end
end
```

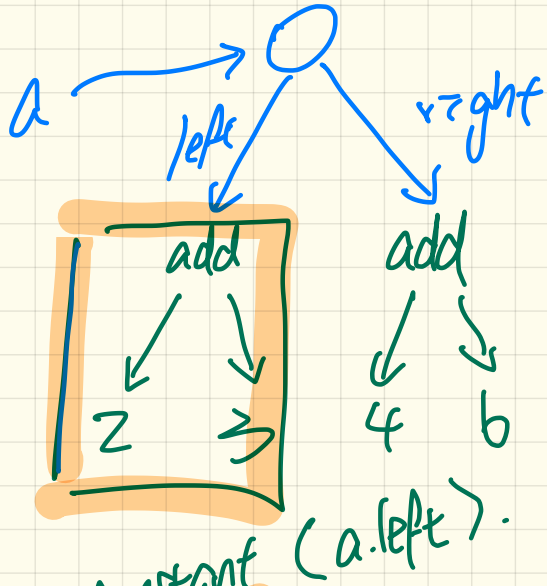
2nd D.S.
visitor
base case no further D.S.
exp.
DT: CONSTANT
2nd D.S.

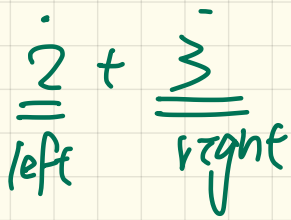
```
class CONSTANT inherit EXPRESSION
...
  accept(v: VISITOR)
    do
      v.visit_constant(Current)
    end
end
```

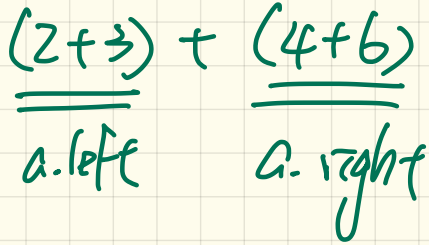
eval_left
DT: EVALUATOR
Alt. v.visit_addition(Cur.)

```
class ADDITION
inherit EXPRESSION COMPOSITE
...
  accept(v: VISITOR)
    do
      v.visit_addition(Current)
    end
end
```

Compilation error.








```
class EVALUATOR inherit VISITOR
  value: INTEGER
  visit_constant(c: CONSTANT)  do value := c.value end
  visit_addition(a: ADDITION)
    local eval_left, eval_right: EVALUATOR
    do a.left.accept(eval_left)
       a.right.accept(eval_right)
       value := eval_left.value + eval_right.value
    end
end
```

VISITOR:
visit_const. (c: CONST.)
Con-EXP:
left: EXP

eval_left.visit_const(a.left)

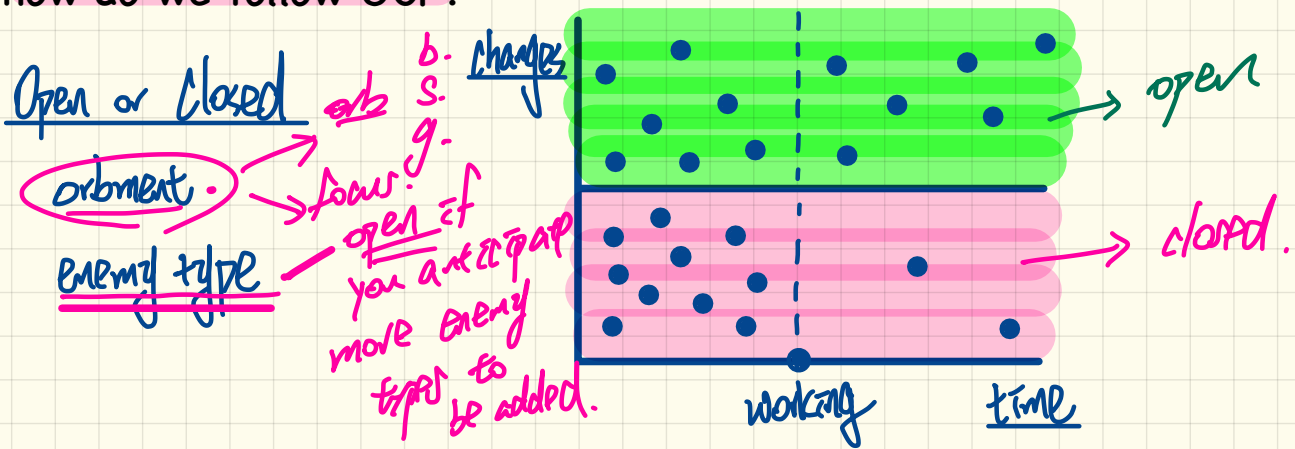
not compile
EXP not descendant of CONST.

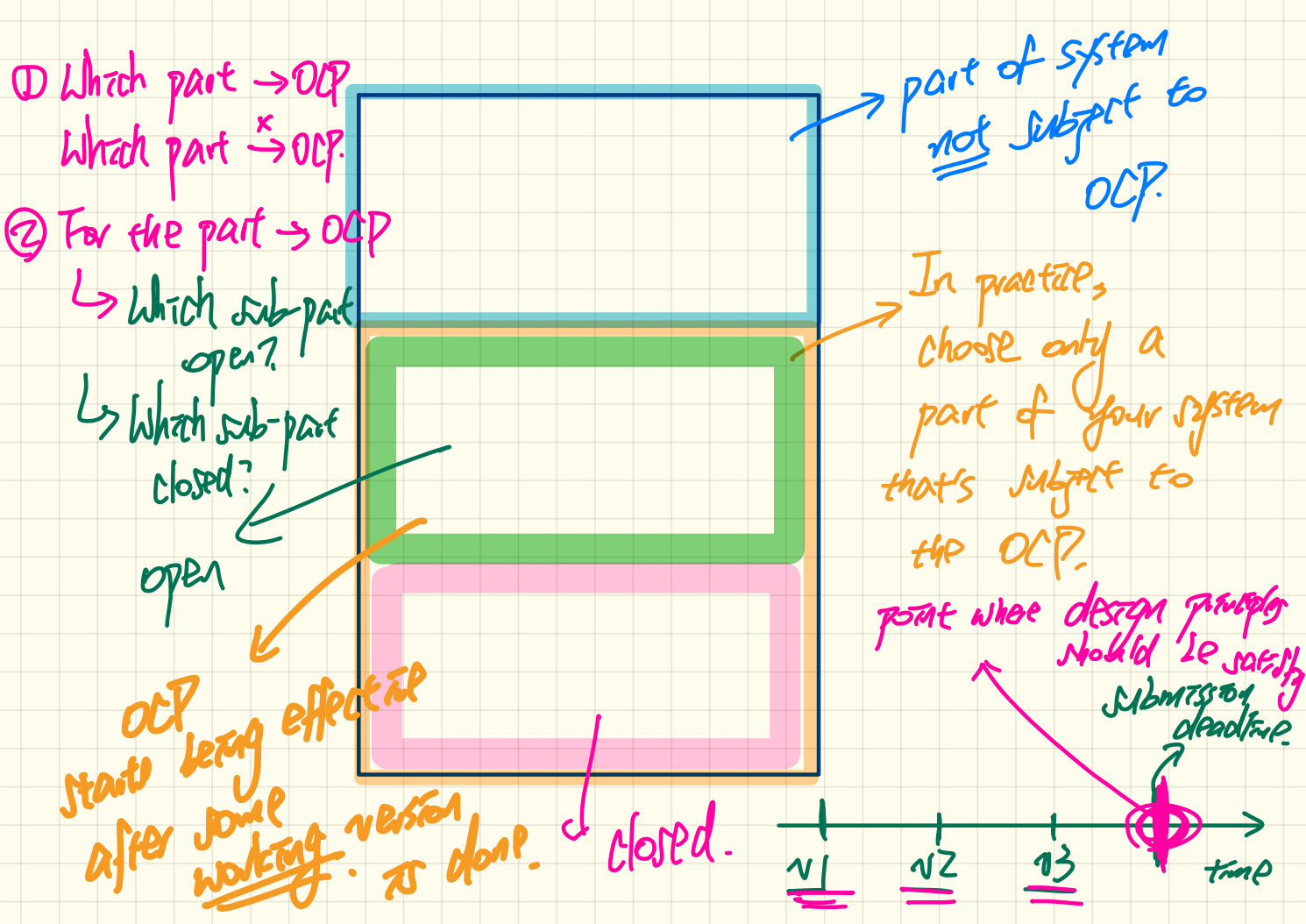
ST: EXP

I think open/close concept is not.. practical to follow.

As an example, in the project (SD2), I am always changing all the parts of all the objects even the very first object I created, because otherwise I cannot satisfy the requirements in a proper way. How do I decide where to draw the line between open and close? What am I missing? On a related question, in a team work project, how do we follow OCP?

in the course of implementing a working version, not yet subject to OCP.




① Which part → OCP
Which part ✗ → OCP.
② For the part → OCP
↳ Which sub-part open?
↳ Which sub-part closed?
open
part of system not subject to OCP.
In practice, choose only a part of your system that's subject to the OCP.
point where design principle should be satisfy
submission deadline
OCP starts being effective after some working version is done.
closed.
v1
v2
v3
time